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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
09/535,765	03/28/2000	Vincent E. Hummel		5136

7590   01/22/2004

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EXAMINER
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ART UNIT	PAPER NUMBER
2183	

DATE MAILED: 01/22/2004

Please find below and/or attached an Office communication concerning this application or proceeding.

PTO-90C (Rev. 10/03)

-- Th MAILING DATE of this communication appears n the cover sheet with the correspondence address --

Peri d f r Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE _3_ MONTH(S) FROM THE MAILING DATE OF THIS COMMUNICATION.
- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If the period for reply specified above is less than thirty (30) days, a reply within the statutory minimum of thirty (30) days will be considered timely.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133).
- Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

1)☒ Responsive to communication(s) filed on _24 November 2003_.

2a)☐ This action is **FINAL.**   2b)☒ This action is non-final.

3)☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under _Ex parte Quayle_, 1935 C.D. 11, 453 O.G. 213.

Disp sition of Claims

4)☒ Claim(s) _1,2,4-9,11-16 and 18-21_ is/are pending in the application.

    4a) Of the above claim(s) _____ is/are withdrawn from consideration.

5)☐ Claim(s) _____ is/are allowed.

6)☒ Claim(s) _1,2,4-9,11-16 and 18-21_ is/are rejected.

7)☐ Claim(s) _____ is/are objected to.

8)☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

9)☐ The specification is objected to by the Examiner.

10)☒ The drawing(s) filed on _28 March 2000_ is/are: a)☐ accepted or b)☒ objected to by the Examiner.

    Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).

    Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).

11)☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. §§ 119 and 120

12)☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).

    a)☐ All b)☐ Some * c)☐ None of:

      1.☐ Certified copies of the priority documents have been received.

      2.☐ Certified copies of the priority documents have been received in Application No. _____.

      3.☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

    * See the attached detailed Office action for a list of the certified copies not received.

13)☐ Acknowledgment is made of a claim for domestic priority under 35 U.S.C. § 119(e) (to a provisional application) since a specific reference was included in the first sentence of the specification or in an Application Data Sheet. 37 CFR 1.78.

    a) ☐ The translation of the foreign language provisional application has been received.

14)☐ Acknowledgment is made of a claim for domestic priority under 35 U.S.C. §§ 120 and/or 121 since a specific reference was included in the first sentence of the specification or in an Application Data Sheet. 37 CFR 1.78.

Attachment(s)

1) ☒ Notice of References Cited (PTO-892)

2) ☐ Notice of Draftsperson's Patent Drawing Review (PTO-948)

3) ☐ Information Disclosure Statement(s) (PTO-1449) Paper No(s) _____ .

4) ☐ Interview Summary (PTO-413) Paper No(s). _____ .

5) ☐ Notice of Informal Patent Application (PTO-152)

6) ☐ Other: .

DETAILED ACTION

1.   Claims 1-2, 4-9, 11-16, 18-21 have been examined.

Papers Submitted

2.   It is hereby acknowledged that the following papers have been received and placed of

record in the file: #8. Request for Continued Examination as received on 10/14/2003 and #9.

Amendment "B" as received on 10/14/2003.

Drawings

3.   The drawings are objected to under 37 CFR 1.83(a). The drawings must show every

feature of the invention specified in the claims. Therefore, the replacement field for a table entry

must be shown or the feature(s) canceled from the claim(s). No new matter should be entered.

A proposed drawing correction or corrected drawings are required in reply to the Office

action to avoid abandonment of the application. The objection to the drawings will not be held

in abeyance.

Claim Rejections - 35 USC § 102

4.   The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the

basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless –

(e) the invention was described in (1) an application for patent, published under section 122(b), by another filed
in the United States before the invention by the applicant for patent or (2) a patent granted on an application for
patent by another filed in the United States before the invention by the applicant for patent, except that an
international application filed under the treaty defined in section 351(a) shall have the effects for purposes of this
subsection of an application filed in the United States only if the international application designated the United
States and was published under Article 21(2) of such treaty in the English language.

5.   Claims 1-2, 6-8, and 20 are rejected under 35 U.S.C. 102(e) as being anticipated by

McFarling, U.S. Patent No. 6,374,349 (as applied in the previous Office Action).

6.   Referring to claim 1, McFarling has taught a method comprising:

a) providing a first taken/not-taken prediction responsive to an address using a saturating counter

branch predictor. See column 9, lines 10-16. It has been disclosed that a bimodal predictor

always provides a prediction. Furthermore, from column 3, lines 11-28, the bimodal predictor is

shown to be an array of saturated counters.

b) providing (1) a second taken/not-taken prediction responsive to the address resulting in a hit in

a local branch history table, and (2) a hit/miss indication for the address. See column 9, lines 18-

24, and Fig.10. Note in Fig.10 that the hit indication controls which prediction is propagated

through the multiplexer.

c) selecting for the address one of (1) the second prediction if the indication is a hit, and (2) the

first prediction if the indication is a miss. Again, see column 9, lines 10-24, and Fig.10. If a

history-table miss occurs, the counter prediction will be used. On the other hand, if a history-

table hit occurs, the history table prediction will be used.

d) updating a field for a matching entry in the local branch history table only if the first

prediction is incorrect, indicating that the entry is used to make a prediction. See column 9, lines

45-52. When the first prediction from the bimodal predictor (saturating counter predictor) is

correct, the corresponding entry in the history table is not replaced. Therefore, replacements

would only occur when the first prediction is incorrect. According to "The American Heritage®

Dictionary of the English Language, Third Edition," 1992, the word "updating" is defined as "to

bring up to date." Therefore, it should be realized that performing a replacement in the local

cache is the same as updating the local cache since the replacement is done to bring the cache up to date. More specifically, when a first branch's prediction information is replaced by a second branch's prediction information, the entry which is storing this prediction information is updated. This update also indicates that this entry's new prediction information will be used in the future to make a prediction for the corresponding branch instruction.

7.   Referring to claim 2, McFarling has taught a method as described in claim 1. McFarling has further taught hashing the address prior to indexing at least one of the saturating counter branch predictor and the local branch history table. See column 3, lines 12-16, and column 4, lines 15-19. Recall that hashing is simply the function of converting a value (in this case, the branch instruction address) into a value used for locating corresponding data in a structure (in this case, the saturating counter predictor and the local history predictor). In the system of McFarling, a portion of a full branch instruction address is used to access a prediction entry in a corresponding structure would be considered hashing.

8.   Referring to claim 6, McFarling has taught a processor comprising:

a) an instruction pointer (IP) generator capable of providing an address. See column 3, lines 12-16. Note that McFarling makes reference to a program counter, which performs the same function as an IP generator (i.e. they both provide an address at which an instruction will be fetched).

b) saturating counter branch prediction (SCBP) logic having an input coupled to the IP generator and capable of providing a first taken/not-taken prediction at an output responsive to the address. See column 9, lines 10-16. It has been disclosed that a bimodal predictor always provides a prediction. Furthermore, from column 3, lines 11-28, the bimodal predictor is shown to be an

array of saturated counters. Furthermore, column 3, lines 12-16, explain that the SCBP logic is

coupled to the program counter (IP generator).

c) local branch history prediction (LBHP) logic having an input coupled to the IP generator and

capable of providing (1) a second taken/not-taken prediction at an output responsive to the

address resulting in a hit, and (2) a hit/miss indication for the address, wherein the LBHP logic

includes at least one local branch history table to provide a taken/not-taken history from a

matching entry in said table in response to a hit. See column 9, lines 18-24, and Fig. 10. Note in

Fig. 10 that the hit indication controls which prediction is propagated through the multiplexer.

Also, in column 4, lines 15-19, McFarling has explained that the local history table is indexed

based on a portion of the branch instruction address (tag). Therefore, the local history table is

also coupled to the IP generator. It should be realized that the LBHP includes at least one local

branch history table to provide a taken/not-taken history in response to a hit. From Fig. 10, it can

be seen that the tag/history data is stored in one table 100 and the prediction (CNT) is stored in a

second table. The most significant bit of CNT is used to make a taken/not-taken prediction.

Finally, it should be realized that the taken/not-taken history is provided from a matching entry

(storage location) in the local table. This matching entry is the location that the branch

instruction's address maps to.

d) a multiplexer having an input coupled to the outputs of the SCBP and LBHP logic and a select

input coupled to receive the hit/miss indication and in response provide (1) the second prediction

if there is a hit and (2) the first prediction if there is a miss. See Fig. 10 and note that the

multiplexer has inputs from both the SCBP (104) and LBHP (104) and that one of the inputs is

chosen based on the hit signal that is used as a select line for the multiplexer. See column 9,

lines 10-24 for further explanation.

e) entry replacement logic to update a replacement field for said matching entry in the at least

one table only if the first prediction is incorrect. See column 9, lines 45-52. When the first

prediction from the bimodal predictor (saturating counter predictor) is correct, the corresponding

entry in the history table is not updated. Therefore, replacements would only occur when the

first prediction is incorrect. In addition, if the matching entry is the least recently used entry,

then this

9.   Referring to claim 7, McFarling has taught a processor as described in claim 6.

McFarling has further taught address hash logic coupled between the IP generator and the inputs

of the SCBP and LBHP logic to provide a plurality of index values to at least one of the SCBP

and LBHP logic. Again, as discussed in column 3, lines 12-16, and column 4, lines 15-19, the

prediction logic is provided an index by performing a truncation hashing function. More

specifically, it has been taught (in the above passages) that the prediction logic receives a certain

amount of low order branch instruction address bits as an index instead of the entire address.

Therefore, the subset of wires used to carry the index from the IP generator to the prediction

logic is address hash logic that performs a truncation algorithm. This logic transforms an initial

value (branch instruction address) into another value used for locating corresponding data in a

structure (SCBP and LBHP).

10.   Referring to claim 8, McFarling has taught a processor as described in claim 6.

McFarling has further taught that the SCBP logic includes a bimodal predictor. See column 9,

lines 10-24.

11.   Referring to claim 20, McFarling has taught a method comprising:

a) providing a first taken/not-taken prediction responsive to an address using a saturating counter branch predictor. See column 9, lines 10-16. It has been disclosed that a bimodal predictor always provides a prediction. Furthermore, from column 3, lines 11-28, the bimodal predictor is shown to be an array of saturated counters.

b) providing a second taken/not-taken prediction responsive to the address resulting in a hit in a local branch history table. See column 9, lines 18-24, and Fig.10. Note in Fig.10 that the hit indication controls which prediction is propagated through the multiplexer.

c) selecting for the address one of the first prediction and the second prediction depending on whether or not the address resulted in said hit. See column 9, lines 10-24, and Fig.10. If a history-table miss occurs, the counter prediction will be used. On the other hand, if a history-table hit occurs, the history table prediction will be used.

d) updating a replacement field, for an existing entry in the local branch history table that includes a history field containing a plurality of taken/not-taken outcomes of one or more previously executed branch instructions associated with said address, only if the first prediction is incorrect. See column 9, lines 45-52. When the first prediction from the bimodal predictor (saturating counter predictor) is correct, the corresponding entry in the history table is not replaced. Therefore, replacements would only occur when the first prediction is incorrect. According to "The American Heritage® Dictionary of the English Language, Third Edition," 1992, the word "updating" is defined as "to bring up to date." Therefore, it should be realized that performing a replacement in the local cache is the same as updating the local cache since the replacement is done to bring the cache up to date. More specifically, when a first branch's

prediction information is replaced by a second branch's prediction information, the entry which

is storing this prediction information is updated. This update also indicates that this entry's new

prediction information will be used in the future to make a prediction for the corresponding

branch instruction.

Claim Rejections - 35 USC § 103

12.   The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all

obviousness rejections set forth in this Office action:

> (a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in
> section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are
> such that the subject matter as a whole would have been obvious at the time the invention was made to a person
> having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the
> manner in which the invention was made.

13.   Claims 4-5 and 11-12 are rejected under 35 U.S.C. 103(a) as being unpatentable over

McFarling, as applied above, in view of Hennessy and Patterson, Computer Architecture - A

Quantitative Approach, 2nd Edition (as applied in the previous Office Action and herein referred

to as Hennessy).

14.   Referring to claim 4, McFarling has taught a method as described in claim 1.

a) Furthermore, it is inherent that an instruction for which a prediction is made would have been

fetched at the corresponding address. This concept is also supported in column 3, lines 12-16.

The bits that are used to index the predictors are taken from an address that is stored in the

program counter (PC). The PC is a register that is used to store a pointer to a location in memory

from which an instruction is fetched.

b) In addition, McFarling has taught that if the instruction is a branch, a determination as to

whether the branch is taken or not-taken will not be available until the instruction has progressed

beyond a decode stage. See column 1, lines 33-36, and note that the outcome of the branch is unknown until after the branch is executed. Since execution inherently follows decoding, the branch determination is not known until the instruction has progressed beyond the decode stage.

c) Finally, although it is inherent that McFarling's system decodes fetched instructions (in order to determine what operation would be performed), McFarling has not explicitly taught that at least one of the first and second predictions is available when the at least one instruction is being decoded. However, Hennessy has shown the state of a pipeline that implements branch prediction. See Figure 3.26 on page 167. Note that during the decode stage (ID) of the branch instruction, the predicted instruction (Instruction i+1) is being fetched in the fetch stage (IF). In order for the predicted instruction to be fetched during the decode stage of the branch instruction, the prediction must be available when the branch is being decoded. It can be seen that the pipeline is kept full and free of stalls through the use of branch prediction. Furthermore, from column 1, lines 49-56, McFarling has explained that the purpose of a branch prediction scheme is to keep the pipeline full and free from stalls, which in fact is supported by the teachings of Hennessy. Therefore, it would have been obvious to one of ordinary skill in the art at the time of the invention to make the at least one of the predictions available when the at least one instruction is being decoded. This will ensure that the pipeline will stay full, as is desired by McFarling. It should be further noted that the scheme taught by Hennessy does not only apply to a system in which branch outcomes are known at the end of the decode stage. This can be realized since McFarling has taught a taken/not-taken prediction scheme in which branch outcomes are unknown until after the branch executes (column 1, lines 33-45). The general idea of such a prediction scheme is to keep the pipeline full regardless of when the branch outcome is

"officially" known, and if the prediction is available while the branch is being decoded, then the

next instruction can be fetched without stalling.

15.   Referring to claim 5, McFarling in view of Hennessy has taught a method as described in

claim 4. If the at least one instruction is a branch, then:

a) it is inherent that the method further comprises determining a target address of the branch. If a

target address were not determined for a conditional branch, then the system would not know

where to fetch the next instruction from when the branch is predicted taken.

b) it is inherent that the method further comprises loading an instruction pointer generator with

the target address if at least one of the first and second predictions indicates that the branch is to

be taken. In order to fetch the target instruction of a predicted-taken branch, the target address

must be loaded into the IP generator (also known as the program counter (PC)). The IP

generator (PC) is a register that is used to store a pointer to a location in memory from which an

instruction is fetched. For the applicant's benefit, please see Hennessy, page 162 and Figure

3.22 on page 163. From the figure it can be seen that the PC (IP generator) is loaded with the

output of a multiplexer. Inputs to the multiplexer include the next instruction address and the

target address (which is outputted by the ADD logic in the ID stage of the pipeline). If the target

address were not stored in the IP generator (PC), then the target instruction could not be fetched

from instruction memory (as shown in Figure 3.22 of Hennessy).

16.   Referring to claim 11, McFarling has taught a processor as described in claim 6.

Furthermore, the processor of claim 11 performs the method described in claim 4. Therefore,

claim 11 is rejected for the same reasons set forth in the rejection of claim 4.

17.   Referring to claim 12, McFarling has taught a processor as described in claim 11.

Furthermore, the processor of claim 12 performs the method described in claim 5. In addition, it

should be realized from the rejection of claims 4 and 5 that the decode stage taught by Hennessy

has the ability to determine a target address. See Figure 3.22 on page 163. Therefore, claim 12

is rejected for the same reasons set forth in the rejection of claim 5.

18.   Claim 9 is rejected under 35 U.S.C. 103(a) as being unpatentable over McFarling, as

applied above, in view of Gochman et al., U.S. Patent No. 5,842,008 (as applied in the previous

office action and herein referred to as Gochman).

19.   Referring to claim 9, McFarling has taught a processor as described in claim 6.

a)   1) McFarling has taught a single branch history table that provides a tag and a taken/not-

taken history associated with the tag in response to a hit. See Fig.10. McFarling has not

explicitly taught the concept of using multiple branch history tables to provide a tag and a

taken/not taken history associated with the tag in response to a hit. However, Gochman has

taught the concept of using multiple memories that are simultaneously accessed based on an

index. After each memory produces its respective data, a single value is selected and

propagated. See Fig.4a and Fig.5. A person of ordinary skill in the art would have recognized

that by using multiple memory tables, as opposed to a single table, the number of index bits used

to access an entry in each bank would be reduced, resulting in a reduction in the amount of

hardware. Therefore, it would have been obvious to one of ordinary skill in the art at the time of

the invention to use multiple branch history tables instead of a single table.

2) In addition, it would have been obvious to one of ordinary skill in the art at the time of

the invention to separate the single history table as taught by McFarling into multiple history

tables since it has been held that making separable where needed is obvious. See Nerwin v.

Erlichman, 168 USPQ 177 (1969). It should be realized that the single-table system serves the

same purpose as the multiple-table system in that an index is applied and a corresponding entry

is selected. Although multiple entries are selected in response to applying an index to multiple

tables (in Applicant's system), the final result is still only a single entry, as taught by McFarling.

b) it is inherent that McFarling's system includes compare logic that is connected to the history

table in order to determine the hit/miss indication. A hit indication is given when the branch

instruction in question has a corresponding entry in the history table. Therefore, in order to give

a hit signal, some logic must be used to compare the entries in the table with the current branch

instruction. Since McFarling has not taught multiple history tables, it follows that he has not

taught multiple compare logic units that are connected to those tables. However, if multiple

tables were used, as discussed in part (a) above, then it would be inherent that multiple compare

logic units would be connected to those tables in order to determine if a branch has a

corresponding entry.

c) McFarling has not taught a history multiplexer coupled to each of the plurality of tables to

provide the history for the hit. However, since McFarling has only taught one table, there is no

need to multiplex multiple values from multiple tables. On the other hand, if multiple tables

were implemented in McFarling's system, as described in part (a) above, then it would follow

that a multiplex system would be needed in order to choose one of the plurality of history values.

This concept has been taught by Gochman in Fig.4a (component 320) and Fig.5. When each of

the values are retrieved from the memories, the selection logic determines which value will be sent through. Therefore, if multiple tables were used in McFarling's system, in order to ensure only the correct history value is used to make the prediction, it would have been obvious to one of ordinary skill in the art at the time of the invention to use a multiplexer to perform the selection.

d) McFarling has taught combinational logic coupled to an output of the history multiplexer to provide the second taken/not-taken prediction. See Fig.10 (prediction line 102 exits from combination logic that produces the taken/not-taken signal). Since McFarling has not taught multiple history tables and a history multiplexer, it follows that he has not taught combinational logic that is connected to a history multiplexer. However, if multiple tables and a history multiplexer were used, as discussed in parts (a) and (c) above, then it would follow that the combinational logic would be coupled to an output of the history multiplexer to provide the second taken/not-taken prediction. For instance, this concept is also shown in Fig.4a, component 330, of Gochman, which controls the prediction based on the output of the selection logic. Therefore, if multiple tables and a history multiplexer were used, it would have been obvious to one of ordinary skill in the art at the time of the invention to connect the combinational logic taught by McFarling to the output of the history multiplexer since the combinational logic uses the history information in order to generate a taken/not-taken prediction.

20.   Claim 13 is rejected under 35 U.S.C. 103(a) as being unpatentable over McFarling, as applied above, in view of Rahman et al., U.S. Patent No. 5,805,878 (as applied in the previous office action and herein referred to as Rahman).

21.   Referring to claim 13, McFarling has taught a processor as described in claim 6.

McFarling has not explicitly taught that the address points to a cache line having a plurality of

instructions. However, Rahman has taught the implementation of a cache with multiple

instructions per cache line. See column 9, lines 17-21. Also, Rahman has suggested in column

9, lines 23-37, that a benefit of such a feature would be to allow for simultaneous fetching of

multiple instructions. This would result in less time spent making memory accesses and

increased instruction-level parallelism, which would result in higher throughput. Therefore, in

order to increase throughput, it would have been obvious to one of ordinary skill in the art at the

time of the invention to implement a cache that contains a plurality of instructions per cache line.

22.   Claims 14-15 and 21 are rejected under 35 U.S.C. 103(a) as being unpatentable over

McFarling, as applied above, in view of Henry et al., U.S. Patent No. 6,550,004 (herein referred

to as Henry).

23.   Referring to claim 14, McFarling has taught an apparatus comprising:

a) means for providing an address of at least one instruction. See column 3, lines 12-16. Note

that McFarling makes reference to a program counter, which performs the same function as an IP

generator (i.e. they both provide an address at which an instruction will be fetched).

b) means for providing a first taken/not-taken branch prediction based upon the current state of a

state machine and responsive to the address. See column 9, lines 10-16. It has been disclosed

that a bimodal predictor always provides a prediction. Furthermore, from column 3, lines 11-28,

the bimodal predictor is shown to be an array of saturated counters. Saturated counters are state

machines in that they function differently when in different states. See Fig.2.

c) local branch history prediction (LBHP) logic having an input coupled to the address providing

means and capable of providing (1) a second taken/not-taken prediction at an output responsive

to the address resulting in a hit, and (2) a hit/miss indication for the address, wherein the LBHP

logic includes at least one local branch history table to provide a taken/not-taken history in

response to a hit. See column 9, lines 18-24, and Fig.10. Note in Fig.10 that the hit indication

controls which prediction is propagated through the multiplexer. Also, in column 4, lines 15-19,

McFarling has explained that the local history table is indexed based on a portion of the branch

instruction address. Therefore, the local history table is also coupled to the IP generator. It

should be realized that the LBHP includes at least one local branch history table to provide a

taken/not-taken history in response to a hit. From Fig.10, it can be seen that the tag/history data

is stored in one table 100 and the prediction (CNT) is stored in a second table. The most

significant bit of CNT is used to make a taken/not-taken prediction.

d) a multiplexer having an input coupled to the outputs of the first prediction means and the

LBHP logic and a select input coupled to receive the hit/miss indication and in response provide

(1) the second prediction if there is a hit and (2) the first prediction if there is a miss. See Fig.10

and note that the multiplexer has inputs from both the SCBP (104) and LBHP (104) and that one

of the inputs is chosen based on the hit signal that is used as a select line for the multiplexer. See

column 9, lines 10-24 for further explanation.

e) McFarling has not explicitly taught means for updating said history field, with the outcome of

an executed branch instruction that is pointed to by said address, only if the first prediction is

incorrect and the second prediction is correct. However, Henry has taught such a concept. See

Fig.6 and note that if the final (first) prediction is incorrect (step 624) and the non-selected

(second) prediction is correct (step 628), then the corresponding history field within predictor

212 (Fig.2) will be updated (step 632). It should be noted from Fig.3 that component 212 does

have history entries. Henry has further disclosed that such an update scheme advantageously

yields improved branch prediction results. As a result, it would have been obvious to one of

ordinary skill in the art at the time of the invention to modify McFarling to employ the update

policy taught by Henry.

24.   Referring to claim 15, McFarling in view of Henry has taught an apparatus as described

in claim 14. Furthermore, claim 15 is rejected for the same reasons set forth in the rejection of

claim 2. Note that hashing and encoding perform the same operation in that a starting value

(branch instruction address) is transformed into another value (index to prediction logic).

25.   Referring to claim 21, McFarling has taught a method comprising:

a) providing a first taken/not-taken prediction responsive to an address using a saturating counter

branch predictor. See column 9, lines 10-16. It has been disclosed that a bimodal predictor

always provides a prediction. Furthermore, from column 3, lines 11-28, the bimodal predictor is

shown to be an array of saturated counters. Saturated counters are state machines in that they

function differently when in different states. See Fig.2.

b) providing (1) a second taken/not-taken prediction responsive to the address resulting in a hit in

a local branch history table, and (2) a hit/miss indication for the address, wherein the second

prediction is based on a taken/not-taken history contained in a history field of an existing entry in

said table. See column 9, lines 18-24, and Fig.10. Note in Fig.10 that the hit indication controls

which prediction is propagated through the multiplexer. Also, in column 4, lines 15-19,

McFarling has explained that the local history table is indexed based on a portion of the branch

instruction address. Therefore, the local history table is also coupled to the IP generator. It

should be realized that the LBHP includes at least one local branch history table to provide a

taken/not-taken history in response to a hit. From Fig.10, it can be seen that the tag/history data

is stored in one table 100 and the prediction (CNT) is stored in a second table. The most

significant bit of CNT is used to make a taken/not-taken prediction.

c) selecting for the address one of (1) the second prediction if the indication is a hit, and (2) the

first prediction if the indication is a miss. Again, see column 9, lines 10-24, and Fig.10. If a

history-table miss occurs, the counter prediction will be used. On the other hand, if a history-

table hit occurs, the history table prediction will be used.

d) McFarling has not explicitly taught means for updating said history field based on an outcome

of an executed branch instruction that is pointed to by said address, only if the first prediction is

incorrect and the second prediction is correct. However, Henry has taught such a concept. See

Fig.6 and note that if the final (first) prediction is incorrect (step 624) and the non-selected

(second) prediction is correct (step 628), then the corresponding history field within predictor

212 (Fig.2) will be updated (step 632). It should be noted from Fig.3 that component 212 does

have history entries. Henry has further disclosed that such an update scheme advantageously

yields improved branch prediction results. As a result, it would have been obvious to one of

ordinary skill in the art at the time of the invention to modify McFarling to employ the update

policy taught by Henry.

26.   Claim 16 is rejected under 35 U.S.C. 103(a) as being unpatentable over McFarling in

view of Henry, as applied above, in view of Gochman, as applied above.

27.   Referring to claim 16, McFarling has taught an apparatus as described in claim 14.

Furthermore, claim 16 is rejected for the same reasons set forth in the rejection of claim 9.

28.   Claims 18-19 are rejected under 35 U.S.C. 103(a) as being unpatentable over McFarling

in view of Henry, as applied above, in view of Hennessy, as applied above.

29.   Referring to claim 18, McFarling has taught an apparatus as described in claim 14.

Furthermore, the apparatus of claim 18 performs the method of claim 4. Therefore, claim 18 is

rejected for the same reasons set forth in the rejection of claim 4.

30.   Referring to claim 19, McFarling has taught an apparatus as described in claim 18.

Furthermore, the apparatus of claim 19 performs the method of claim 5. Therefore, claim 19 is

rejected for the same reasons set forth in the rejection of claim 5.

Conclusion

31.   The prior art made of record and not relied upon is considered pertinent to applicant's

disclosure. Applicant is reminded that in amending in response to a rejection of claims, the

patentable novelty must be clearly shown in view of the state of the art disclosed by the

references cited and the objections made. Applicant must also show how the amendments avoid

such references and objections. See 37 CFR § 1.111(c).

        Driesen et al., The Cascaded Predictor: Economical and Adaptive Branch Target

Prediction, December 1998, pp.1-10, has taught a cascaded predictor wherein a second stage

predictor provides a prediction if it is able to; otherwise, the first-stage predictor will provide a
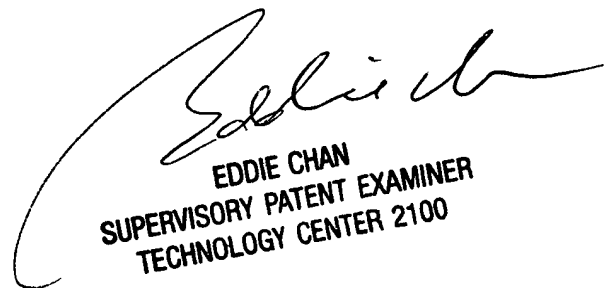
prediction.

Rappoport et al., U.S. Patent No. 6,601,161, has taught a method and system for branch target prediction using path information.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to David J. Huisman whose telephone number is (703) 305-7811. The examiner can normally be reached on Monday-Friday (8:00-4:30).

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Eddie Chan can be reached on (703) 305-9712. The fax phone number for the organization where this application or proceeding is assigned is (703) 746-7239.

Any inquiry of a general nature or relating to the status of this application or proceeding should be directed to the receptionist whose telephone number is (703) 305-3900.

DJH
David J. Huisman
January 6, 2004

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